

RATIONAL INTERPOLATION TO $|x|$ AT THE CHEBYSHEV NODES

LEV BRUTMAN AND ELI PASSOW

Recently the authors considered Newman-type rational interpolation to $|x|$ induced by arbitrary sets of interpolation nodes and showed that under mild restrictions on the location of the interpolation nodes, the corresponding sequence of rational interpolants converges to $|x|$. In the present paper we consider the special case of the Chebyshev nodes which are known to be very efficient for polynomial interpolation. It is shown that, in contrast to the polynomial case, the approximation of $|x|$ induced by rational interpolation at the Chebyshev nodes has the same order as rational interpolation at equidistant points.

1. INTRODUCTION

The function $|x|$ has been the focus of much research in approximation theory over the years. Its fundamental role in polynomial approximation is well illustrated by Lebesgue's proof of the Weierstrass approximation theorem, which is based solely on the fact that the single function $|x|$ can be approximated. However, as was shown by Bernstein [1], the order of the best uniform approximation of $|x|$ by *polynomials* is only $O(n^{-1})$.

In contrast to this, Newman [5] demonstrated that *rational* approximation to $|x|$ is much more favorable, namely $|x|$ may be approximated uniformly by rational functions at an exponential rate. Newman's result generated a great deal of research, much of which focused on the problem of sharpening the asymptotic results for the error in the best rational approximation. The most recent result in this direction is the proof of the so-called "8" conjecture by Stahl. (See [7], where the main result is presented and an extensive historical review is given).

In a recent paper [2] the authors considered Newman-type rational approximation induced by arbitrary sets of interpolation points. Let $X = \{0 < x_1^{(n)} < x_2^{(n)} < \dots < x_n^{(n)} \leq 1\}$ be a set of n distinct points in $(0, 1]$ and let $p(x) = \prod_{k=1}^n (x + x_k^{(n)})$. (In the sequel, when there is no possibility for confusion, the superscript (n) will be omitted.) The rational function, corresponding to the set X , is defined by

$$r_n(X; x) = x \frac{p(x) - p(-x)}{p(x) + p(-x)}.$$

Received 28th August, 1996

Copyright Clearance Centre, Inc. Serial-fee code: 0004-9729/97 \$A2.00+0.00.

It can be easily verified that $r_n(X; x)$ interpolates $|x|$ at the following set of $2n + 1$ points: $\{-x_n, \dots, -x_1, 0, x_1, \dots, x_n\}$. Since $r_n(X; x)$ as well as $|x|$ are even functions, the study of the approximation error $e_n(X; x) = |x| - r_n(X; x)$ may be restricted to the interval $[0, 1]$, where it can be represented in the following form:

$$(1) \quad e_n(X; x) = \frac{2xh_n(X; x)}{1 + h_n(X; x)}, \quad 0 \leq x \leq 1,$$

where
$$h_n(X; x) = \frac{p(-x)}{p(x)} = \prod_{k=1}^n \frac{x_k - x}{x_k + x}.$$

In the sequel we shall use the following general estimates which were proved in [2]:

STATEMENT 1.1. Let $S_1 = S_1^{(n)}(X) = \sum_{k=1}^n x_k^{(n)}$. Then

$$(2) \quad |h_n(X; x)| \leq e^{-xS_1}, \quad 0 \leq x \leq 1,$$

$$(3) \quad |e_n(X; x)| \leq \frac{2}{S_1}, \quad -1 \leq x \leq 1.$$

STATEMENT 1.2. Let $A_n = A_n(X) = 1 / \sum_{k=1}^n x_k^{-1}$. Then

$$(4) \quad |e_n(X; x)| \leq 1/A_n, \quad x \in [-x_1, x_1].$$

Note that (3) implies in particular that for the set E of equally spaced points

$$(5) \quad |e_n(E; x)| \leq \frac{4}{n + 1},$$

and thus the function $|x|$ may be uniformly approximated by rational interpolation at the equidistant points with the rate at least $O(1/n)$. This is in striking contrast to the classical result of Bernstein that the sequence of Lagrange interpolating polynomials to $|x|$ at equally spaced points in $[-1, 1]$ diverges everywhere, except at zero and the end-points. (See for example, [3].) It should be mentioned that estimate (5) is a bit conservative, since it was proved by Werner in [8] that the exact order of rational interpolation of $|x|$ at equidistant points is $O(1/n \log n)$.

In the present paper we consider rational interpolation to $|x|$ corresponding to the set of the Chebyshev nodes which are known to be very efficient for polynomial interpolation. We show that the exact order of approximation of $|x|$ by rational interpolation at the Chebyshev nodes is also $O(1/n \log n)$. Thus, in contrast to the polynomial case, for rational interpolation of $|x|$ the Chebyshev nodes are not better than the equidistant ones.

Finally we would like to mention that the method of our proof is rather general and may be applied to other specific sets of interpolation points.

2. RESULTS

Consider the case of the Chebyshev nodes, namely let

$$X = T := \{x_k\} = \sin((2k - 1)\pi/(4n)), \quad k = 1, 2, \dots, n$$

be the roots of the Chebyshev polynomial $T_{2n}(x)$ of degree $2n$, lying in $(0, 1)$. Then, as can be easily verified,

$$p(x)p(-x) = \frac{(-1)^n T_{2n}(x)}{2^{2n-1}},$$

and therefore we have

$$(6) \quad h_n(T; x) = \frac{(-1)^n T_{2n}(x)}{2^{2n-1} p^2(x)}.$$

Since $|x|$ and $r_n(X; x)$ are even functions in $[-1, 1]$ we can restrict ourselves to $x \in [0, 1]$. The following estimate holds:

LEMMA 2.1.

$$(7) \quad |h_n(T; x)| \leq \frac{1}{4n}, \quad x \in [x_1, 1].$$

PROOF: Since $|T_n(x)| \leq 1$ for $x \in [-1, 1]$, it follows from (6) that for $x \geq x_1$

$$(8) \quad |h_n(T; x)| \leq \frac{1}{2^{2n-1} p^2(x)} = \frac{1}{2^{2n-1} \prod_{k=1}^n (x + x_k)^2} \leq \frac{1}{2^{2n-1} \prod_{k=1}^n (x_1 + x_k)^2}.$$

Let $B_n := \prod_{k=1}^n (x_1 + x_k)$. An easy computation reveals:

$$\begin{aligned} B_n &= \prod_{k=1}^n \left[\sin \frac{\pi}{4n} + \sin \frac{(2k - 1)\pi}{4n} \right] = \prod_{k=1}^n 2 \sin \frac{k\pi}{4n} \prod_{k=2}^n \cos \frac{(k - 1)\pi}{4n} \\ &= \sqrt{2} \prod_{k=1}^{n-1} 2 \sin \frac{k\pi}{4n} \cos \frac{k\pi}{4n} = \frac{\sqrt{2}}{2} \prod_{k=1}^{2n-1} \sin \frac{k\pi}{2n}. \end{aligned}$$

By using the well-known identity (see for example, formula (1.392) in [4])

$$\prod_{k=1}^{2n-1} \sin \frac{k\pi}{2n} = \frac{n}{2^{2n-2}},$$

we find

$$(9) \quad B_n^2 = \frac{n}{2^{2n-3}}.$$

Combining (8) and (9) completes the proof of the lemma. □

Now we are in a position to prove the following estimate for the approximation error:

THEOREM 2.2.

$$(10) \quad |e_n(T; x)| \leq \begin{cases} \frac{C_1}{n \log n}, & |x| \in \left[0, \sin \frac{\pi}{4n}\right], \\ \frac{C_2 \log n}{n^2}, & |x| \in \left[\sin \frac{\pi}{4n}, \frac{\pi \log n}{n}\right], \\ \frac{C_3}{n^2}, & |x| \in \left[\frac{\pi \log n}{n}, 1\right]. \end{cases}$$

PROOF: As before we can restrict our analysis to $x \geq 0$. Consider first the case $x \in [0, x_1] = [0, \sin(\pi/(4n))]$. In order to apply (4), we have to estimate the following sum:

$$(11) \quad \begin{aligned} A_n(T) &:= \sum_{k=1}^n x_k^{-1} = \sum_{k=1}^n \frac{1}{\sin((2k-1)\pi/4n)} = \frac{1}{2} \sum_{k=1}^{2n} \frac{1}{\sin((2k-1)\pi/4n)} \\ &= n \left[\frac{1}{2n} \sum_{k=1}^{2n} \frac{1}{\sin((2k-1)\pi/4n)} \right] = n\Lambda_{2n-1}(T). \end{aligned}$$

Here $\Lambda_{2n-1}(T)$ is the Lebesgue constant for polynomial interpolation corresponding to the set of the Chebyshev nodes, for which the following two-sided inequality holds (see for example, [6]):

$$(12) \quad a_0 + \frac{2}{\pi} \log 2n < \Lambda_{2n-1}(T) \leq 1 + \frac{2}{\pi} \log 2n,$$

where $a_0 = \frac{2}{\pi} \left(\gamma + \log \frac{8}{\pi} \right) = 0.9625 \dots$, γ being Euler's constant.

Thus
$$|e_n(T; x)| \leq \frac{1}{A_n(T)} < \frac{1}{n(a_0 + (2/\pi) \log 2n)} \leq \frac{C_1}{n \log n}.$$

Further consider the case $x > x_1$. Note first that in view of the lemma

$$|e_n(T; x)| \leq \frac{8}{3} |h_n(T; x)|, \quad x \in [x_1, 1],$$

and therefore we can restrict ourselves to finding an upper bound for $|h_n(T; x)|$. To this end we apply (2) and take into account that for the Chebyshev nodes

$$(13) \quad S_1(T) = \sum_{k=1}^n \sin(2k-1) \frac{\pi}{4n} = \sum_{k=1}^n \cos(2k-1) \frac{\pi}{4n} = \frac{1}{2 \sin(\pi/4n)} \geq \frac{2n}{\pi}.$$

In (13) we have used the well-known formula

$$\sum_{k=1}^n \cos(2k-1)t = \frac{\sin 2nt}{2 \sin t}.$$

Thus we obtain

$$|h_n(T; x)| \leq e^{-(2nx/\pi)}, \quad x \in [0, 1].$$

Now we can require

$$e^{-(2nx/\pi)} < \frac{1}{n^2} = e^{-2 \log n},$$

which will be satisfied assuming

$$x > \frac{\pi \log n}{n}.$$

It remains to consider the intermediate interval $x \in [\sin(\pi/4n), (\pi \log n)/n]$. But in this case it follows from (1), in view of the lemma, that

$$|e_n(T; x)| \leq \frac{|2xh_n(T; x)|}{1 - |h_n(T; x)|} \leq \left(\frac{2\pi \log n}{n}\right) \left(\frac{1}{4n}\right) \left(\frac{4}{3}\right) = \frac{2\pi \log n}{3n^2}.$$

This completes the proof of the theorem. □

COROLLARY 2.3. *For any $x \in [-1, 1]$ the following estimate holds:*

$$(14) \quad |e_n(T; x)| \leq \frac{C}{n \log n}.$$

Finally we show that the estimate (14) is sharp, namely, the following result holds:

THEOREM 2.4. *Let $x^* = 1/(n \log n)$. Then*

$$(15) \quad |e_n(T; x^*)| \geq \frac{C}{n \log n}, \quad n \geq n_0.$$

PROOF: Note first that for $n > 3$, $x^* \in [0, x_1]$ and since in this interval $0 \leq h_n(T; x) \leq 1$, we can write:

$$n \log n |e_n(T; x^*)| = \frac{e_n(T; x^*)}{x^*} = \frac{2h_n(T; x^*)}{1 + h_n(T; x^*)} \geq h_n(T; x^*).$$

Thus in order to prove (15) we have to show that the sequence $\{F_n\}_{n=1}^\infty$ defined by

$$(16) \quad F_n := \frac{1}{h_n(T; x^*)} = \frac{p^2(x^*)2^{2n-1}}{(-1)^n T_{2n}(x^*)}$$

is bounded. To this end note that

$$\lim_{n \rightarrow \infty} [(-1)^n T_{2n}(x^*)] = 1,$$

and therefore it suffices to consider the behaviour of the numerator of (16). Let

$$(17) \quad P_n := p^2(x^*) = \prod_{k=1}^n (x^* + x_k)^2 = \prod_{k=1}^n x_k^2 \prod_{k=1}^n \left(1 + \frac{x^*}{x_k}\right)^2 := Q_n * R_n.$$

Then for the first factor we have

$$(18) \quad Q_n = \prod_{k=1}^n \sin^2 \frac{(2k-1)\pi}{4n} = \prod_{k=1}^{2n} \sin \frac{(2k-1)\pi}{4n} = \frac{1}{2^{2n-1}},$$

where the last equality is a well-known formula (see for example, formula 1.392(2) from [4]).

Thus it remains to verify that the sequence $\{R_n\}_{n=1}^\infty$ is bounded. Taking into account (11) and (12), we obtain

$$(19) \quad \begin{aligned} \log R_n &= 2 \sum_{k=1}^n \log \left(1 + \frac{x^*}{x_k}\right) \leq 2x^* \sum_{k=1}^n \frac{1}{x_k} = 2x^* A_n(T) \\ &= 2x^* n \Lambda_{2n-1}(T) \leq \frac{2}{\log n} \left(1 + \frac{2}{\pi} \log 2n\right) \leq C, \end{aligned}$$

and the result follows. □

REFERENCES

- [1] S. Bernstein, ‘Sur la meilleure approximation de $|x|$ par des polynômes de degrés donnés’, *Acta Math.* **37** (1913), 1–57.
- [2] L. Brutman and E. Passow, ‘On rational interpolation to $|x|$ ’, *Constr. Approx.* (to appear).
- [3] G.J. Byrne, T.M. Mills and S.J. Smith, ‘On Lagrange’s interpolation with equidistant nodes’, *Bull. Austral. Math. Soc.* **42** (1990), 81–89.
- [4] I.S. Gradshteyn and I.M. Ryzhik, *Table of integrals, series and products* (Academic Press, New York, 1980).
- [5] D. Newman, ‘Rational approximation to $|x|$ ’, *Michigan Math. J.* **11** (1964), 11–14.
- [6] T.J. Rivlin, *Chebyshev polynomials*, (2nd ed.) (Wiley, New York, 1990).
- [7] H. Stahl, ‘Best uniform rational approximation of $|x|$ on $[-1, 1]$ ’, *Mat. Sb.* **183** (1992), 85–118.
- [8] H. Werner, ‘Rationale Interpolation von $|x|$ in äquidistanten Punkten’, *Math. Z.* **180** (1982), 85–118.

Department of Mathematics and
 Computer Science
 University of Haifa
 Haifa 31905
 Israel
 e-mail: lev@mathcs.haifa.ac.il

Department of Mathematics
 Temple University
 Philadelphia PA 19122
 United States of America
 e-mail: passow@euclid.math.temple.edu