


PAPER

Discrete equational theories

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Abstract

On a locally λ -presentable symmetric monoidal closed category \mathcal{V} , λ -ary enriched equational theories correspond to enriched monads preserving λ -filtered colimits. We introduce discrete λ -ary enriched equational theories where operations are induced by those having discrete arities (equations are not required to have discrete arities) and show that they correspond to enriched monads preserving λ -filtered colimits and surjections. Using it, we prove enriched Birkhoff-type theorems for categories of algebras of discrete theories. This extends known results from metric spaces and posets to general symmetric monoidal closed categories.

Keywords: Enriched equational theory; enriched monad; Birkhoff subcategory

1. Introduction

Motivated by probabilistic programming, Mardare et al. (2016, 2017) introduced quantitative equations and developed universal algebra over metric spaces. In the recent paper Rosický (2021), we have related their quantitative equational theories to the theories of Bourke and Garner (2019). More papers were influenced by Mardare et al. (2016, 2017), in particular Adámek (2022) and Milius and Urbat (2019). In the first one, Adámek showed that ω_1 -basic quantitative equational theories correspond to \aleph_1 -ary enriched monads on metric spaces preserving surjections. In Milius and Urbat (2019), Milius and Urbat clarified the Birkhoff-type theorems stated in Mardare et al. (2017).

In this paper, we work in a general symmetric monoidal closed category \mathcal{V} which is locally λ -presentable as a closed category. The leading examples are **Pos** (posets) and **Met** (metric spaces). Under discrete objects of \mathcal{V} , we mean copowers of the monoidal unit I . In **Pos**, these are discrete posets and in **Met** discrete spaces (having all distances 0 or ∞). We introduce discrete theories whose operations are induced by those with a discrete arity. We show that discrete theories correspond to λ -ary enriched monads on \mathcal{V} preserving surjections (i.e., morphisms for which I is projective). As a special case, we get the result of Adámek (2022). Under mild assumptions, surjections form a left part of a factorization system $(\text{Surj}, \text{Inj})$ on the underlying category \mathcal{V}_0 of \mathcal{V} . We assume that \mathcal{V}_0 is Inj -locally μ -generated for $\mu \leq \lambda$ in the sense of Di Liberti and Rosický (2022) and prove Birkhoff's type theorems for categories of algebras of discrete theories. Metric spaces are locally \aleph_1 -presentable and also Inj -locally \aleph_0 -generated and our Birkhoff-type theorems yield those of Mardare et al. (2017).

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The Appendix withdraws some claims of Rosický (2021, Section 5).

2. Background on Algebraic Theories

Classical universal algebra starts with a *signature* Σ giving a set of operations equipped with arities which are finite cardinals. A Σ -algebra A assigns to every n -ary operation f a mapping $f_A : A^n \rightarrow A$. Then one defines *terms* and *equations*, interprets terms on an algebra and says when an equation is satisfied by an algebra. An *equational theory* E is a set of equations and an algebra satisfies E if it satisfies all equations from E . The forgetful functor $U : \mathbf{Alg}(E) \rightarrow \mathbf{Set}$ from the category $\mathbf{Alg}(E)$ of algebras satisfying E has a left adjoint F and $\mathbf{Alg}(E)$ is equivalent to the category $\mathbf{Alg}(T)$ of algebras for the monad $T = UF$. The algebra Fn consists of equivalence classes of n -ary terms and morphisms $Fm \rightarrow Fn$ give m -tuples of n -ary terms. We can consider them as (n, m) -ary operations where the *input arity* n is the usual arity and the *output arity* m is the multiplicity. Then the superposition of terms is replaced by the composition of these operations. If \mathcal{N} is the full subcategory of \mathbf{Set} consisting of finite cardinals and \mathcal{T} the full subcategory of $\mathbf{Alg}(E)$ consisting of free algebras on finite cardinals, then the domain-codomain restriction of F gives an identity-on-objects functor $J : \mathcal{N} \rightarrow \mathcal{T}$. The dual \mathcal{T}^{op} of \mathcal{T} is a Lawvere theory whose algebras A are functors $\hat{A} : \mathcal{T}^{\text{op}} \rightarrow \mathbf{Set}$ preserving finite products. This is the same as $\hat{A}J^{\text{op}} = \mathbf{Set}(K-, A)$ where $K : \mathcal{N} \rightarrow \mathbf{Set}$ is the inclusion and $A = \hat{A}(1)$. Then $\mathbf{Alg}(E)$ is equivalent to the category of algebras of \mathcal{T}^{op} . Hence, classical universal algebra can be equivalently captured by equational theories, finitary monads (= monads preserving filtered colimits), or Lawvere theories. All this is well known due to Birkhoff, Lawvere and Linton (see Adámek et al. (2011)).

In ordered universal algebra, given a signature Σ , a Σ -algebra is a poset A equipped with monotone mappings $f_A : A^n \rightarrow A$ for every n -ary operation f from Σ . But, instead of equations, one takes inequations $p \leq q$ of terms (see Bloom (1976)). An algebra satisfies this inequation if $p_A \leq q_A$ in the poset of monotone mappings $A^n \rightarrow A$. It is natural to take enriched signatures where arities are finite posets. Then the resulting enriched inequational theories correspond to finitary enriched monads on \mathbf{Pos} (Adámek et al., 2021). Recall that the category \mathbf{Pos} of posets and monotone mapping is locally finitely presentable as a cartesian closed category. The free algebra FX on a finite poset X consists of equivalence classes of X -ary terms and morphisms $FY \rightarrow FX$ can be taken as (X, Y) -ary operations. The inequation $p \leq q$ of X -ary terms, which is called an *inequation in the context* X , then means that the pair p, q of terms is an $(X, 2)$ -ary operation where 2 is a two-element chain. In this way, inequational theories can be replaced by equational theories in (X, Y) -ary operations (see Rosický (2021, Example 4.11(1))). If \mathcal{F} is the full subcategory of \mathbf{Pos} consisting of finite poset and \mathcal{T} the full subcategory of $\mathbf{Alg}(E)$ consisting of free algebras on finite posets then the domain-codomain restriction of F gives an identity-on-objects enriched functor $J : \mathcal{F} \rightarrow \mathcal{T}$. The dual \mathcal{T}^{op} of \mathcal{T} is an enriched Lawvere theory (Power, 2005). Hence, ordered universal algebra can be equivalently captured by inequational theories, enriched equational theories, finitary monads, or enriched Lawvere theories.

In metric universal algebra, one has to take into account that the category \mathbf{Met} of metric spaces (distances ∞ are allowed) and nonexpanding maps is only locally \aleph_1 -presentable as a symmetric monoidal closed category. Thus, one has to take countable metric spaces as arities. Given such an enriched signature Σ , a Σ -algebra is a metric space A equipped with nonexpanding mappings $f_A : A^X \rightarrow A$ for every X -ary operation f from Σ . Now, instead of equations, one takes quantitative equations $p =_\varepsilon q$ of terms where $\varepsilon \geq 0$ is a real number (see Mardare et al. (2016, 2017)). An algebra satisfies this quantitative equation if $d(p_A, q_A) \leq \varepsilon$ in the metric space of nonexpanding mappings $A^X \rightarrow A$. These equations are called *basic quantitative equations* (in context X). This means that the pair p, q of terms is an $(X, 2_\varepsilon)$ -ary operation where 2_ε is a two-element metric space with the distance ε between the two points. In this way, basic quantitative theories can be replaced by equational theories in (X, Y) -ary operations (see Rosický (2021, Example 4.11(2))). If \mathcal{C} is the

full subcategory of **Met** consisting of countable metric spaces and \mathcal{T} the full subcategory of $\mathbf{Alg}(E)$ consisting of free algebras on countable metric spaces then the domain-codomain restriction of F gives an identity-on-objects enriched functor $J : \mathcal{F} \rightarrow \mathcal{T}$. The dual \mathcal{T}^{op} of \mathcal{T} is an enriched \aleph_1 -ary enriched Lawvere theory, and these theories correspond to enriched monads preserving \aleph_1 -filtered colimits (see Power (2005)). Hence, metric universal algebra can be equivalently captured by basic quantitative theories, enriched equational theories, \aleph_1 -ary monads, or enriched \aleph_1 -ary Lawvere theories.

In general, let \mathcal{V} be a symmetric monoidal closed category with the unit object I and the underlying category \mathcal{V}_0 . We assume that \mathcal{V} is locally λ -presentable as a symmetric monoidal closed category which means that the underlying category \mathcal{V}_0 is locally λ -presentable, the tensor unit I is λ -presentable and $X \otimes Y$ is λ -presentable whenever X and Y are λ -presentable. We will denote by \mathcal{V}_λ the (representative) small, full subcategory consisting of λ -presentable objects.

Following Bourke and Garner (2019), let \mathcal{A} be a small, full, dense sub- \mathcal{V} -category of \mathcal{V} with the inclusion $K : \mathcal{A} \rightarrow \mathcal{V}$. Objects of \mathcal{A} are called *arities*. Then an \mathcal{A} -pretheory is an identity-on-objects \mathcal{V} -functor $J : \mathcal{A} \rightarrow \mathcal{T}$. A \mathcal{T} -algebra is an object A of \mathcal{V} together with a \mathcal{V} -functor $\hat{A} : \mathcal{T}^{\text{op}} \rightarrow \mathcal{V}$ whose composition with J^{op} is $\mathcal{V}(K-, A)$ (in Bourke and Garner (2019), \mathcal{T} -algebras are called concrete \mathcal{T} -models). Every \mathcal{A} -pretheory induces a \mathcal{V} -monad $T : \mathcal{V} \rightarrow \mathcal{V}$ given by its \mathcal{V} -category $\mathbf{Alg}(\mathcal{T})$ of algebras. Conversely, a \mathcal{V} -monad T induces an \mathcal{A} -pretheory $J : \mathcal{A} \rightarrow \mathcal{T}$ where \mathcal{T} is the full subcategory of $\mathbf{Alg}(T)$ consisting of free algebras on objects from \mathcal{A} and J is the domain-codomain restriction of the free algebra functor. An \mathcal{A} -pretheory is an \mathcal{A} -theory if it is given by its monad. Then \hat{A} is the hom-functor $\mathbf{Alg}(\mathcal{T})(-, A)$ restricted on free algebras over \mathcal{A} . Conversely, a monad T is \mathcal{A} -nervous if it is given by its \mathcal{A} -theory. In this way, we get a one-to-one correspondence between \mathcal{A} -theories and \mathcal{A} -nervous monads (see Bourke and Garner (2019, Corollary 21)).

Under a λ -ary \mathcal{V} -theory we will mean a \mathcal{V}_λ -theory. Following Bourke and Garner (2019), λ -ary \mathcal{V} -theories correspond to \mathcal{V} -monads on \mathcal{V} preserving λ -filtered colimits. They are called λ -ary \mathcal{V} -monads. Hence, λ -ary monads are precisely \mathcal{V}_λ -nervous monads.

3. Surjections

The underlying functor $\mathcal{V}_0(I, -) : \mathcal{V}_0 \rightarrow \mathbf{Set}$ has a left adjoint $- \cdot I$ sending a set X to the coproduct $X \cdot I$ of X copies of I in \mathcal{V}_0 . Objects $X \cdot I$ will be called *discrete*. Every object V of \mathcal{V} determines a discrete object $V_0 = \mathcal{V}_0(I, V) \cdot I$ and morphism $\delta_V : V_0 \rightarrow V$ given by the counit of the adjunction. Every morphism $f : V \rightarrow W$ determines the morphism $f_0 = \mathcal{V}_0(I, f) \cdot I$ between the underlying discrete objects.

A morphism $f : A \rightarrow B$ will be called a *surjection* if $\mathcal{V}_0(I, f)$ is surjective. Let *Surj* denote the class of all surjections in \mathcal{V}_0 and let *Inj* be the class of morphisms of \mathcal{V}_0 having the unique right lifting property with respect to every surjection. Morphisms from *Inj* will be called *injections*. Recall that g is an injection iff for every surjection f and every commutative square

$$\begin{array}{ccc}
 A & \xrightarrow{u} & C \\
 f \downarrow & & \downarrow g \\
 B & \xrightarrow{v} & D
 \end{array}$$

there is a unique $t : B \rightarrow C$ such that $tf = u$ and $gt = v$.

Lemma 3.1. *Surj is accessible and closed under λ -directed colimits in $\mathcal{V}_0^\rightarrow$.*

Proof. f is a surjection iff it has the right lifting property with respect to $0 \rightarrow I$. Like in Rosický (2009, Proposition 3.3), the result follows from Adámek and Rosický (1994, Proposition 4.7). \square

Remark 3.2. (1) Surjections are right-cancellable, i.e., if gf is a surjection then g is a surjection.
 (2) Surjections are closed under products and stable under pullbacks (because they are given by the right lifting property with respect to $0 \rightarrow I$).

Lemma 3.3. $(Surj, Inj)$ is a factorization system in \mathcal{V}_0 if and only if $Surj$ is closed under colimits in $\mathcal{V}_0^{\rightarrow}$.

Proof. Necessity is evident. Conversely, following 3.1, $Surj$ has a small dense subcategory \mathcal{S} . If $Surj$ is closed under colimits in $\mathcal{V}_0^{\rightarrow}$, then $Surj$ is the closure under \mathcal{S} under colimits in $\mathcal{V}_0^{\rightarrow}$. Following Fajstrup and Rosický (2008, Theorem 2.2), $(Surj, Inj)$ is a factorization system. \square

Lemma 3.4. Injections are monomorphisms.

Proof. Take an injection $g : C \rightarrow D$ and $u, v : B \rightarrow C$ such that $gu = gv$. Consider a commutative square

$$\begin{array}{ccc}
 B + B & \xrightarrow{(u,v)} & C \\
 \nabla \downarrow & & \downarrow g \\
 B & \xrightarrow{gu} & D
 \end{array}$$

where ∇ is the codiagonal. Since ∇ is a split epimorphism, it is a surjection and thus there is a unique diagonal $B \rightarrow C$. Hence, $u = v$. We have proved that g is a monomorphism. \square

Lemma 3.5. Assume that $(Surj, Inj)$ is a factorization system. Then

- (1) $Surj$ contains all strong epimorphisms, and
- (2) if I is a generator in \mathcal{V}_0 then Inj contains all strong monomorphisms.

Proof. In every locally presentable category, (strong epi, mono) and (epi, strong mono) are factorization systems (see Adámek and Rosický (1994, Proposition 1.61)). Thus, (1) follows from 3.4. For (2), it suffices to show that every surjection is an epimorphism, which follows from I being a generator. \square

Recall that I is connected iff $\mathcal{V}_0(I, -)$ preserves coproducts.

Lemma 3.6. If I is connected and every regular epimorphism is a surjection, then $(Surj, Inj)$ is a factorization system.

Proof. If I is connected, then $Surj$ is closed under coproducts in $\mathcal{V}_0^{\rightarrow}$. If regular epimorphisms are surjections, then surjections are closed under coequalizers in $\mathcal{V}_0^{\rightarrow}$. Indeed, let f_0 and f_1 be surjections and

$$f_0 \begin{array}{c} \xrightarrow{(u_0, v_0)} \\ \rightrightarrows \\ \xrightarrow{(u_1, v_1)} \end{array} f_1 \xrightarrow{(u, v)} f$$

be a coequalizer in $\mathcal{V}_0^{\rightarrow}$. Then

$$B_0 \begin{array}{c} \xrightarrow{v_0} \\ \rightrightarrows \\ \xrightarrow{v_1} \end{array} B_1 \xrightarrow{v} B$$

is a coequalizer in \mathcal{V}_0 , hence v is a surjection. Following 3.2(1), f is a surjection. \square

The factorization system $(Surj, Inj)$ is λ -convenient in the sense of Di Liberti and Rosický (2022) if

- (1) \mathcal{V} is *Surj*-cowellpowered, i.e., every object of \mathcal{V} has only a set of surjective quotients, and
- (2) *Inj* is closed under λ -directed colimits, i.e., every λ -directed colimit of injections has the property that a colimit cocone
 - (a) consists of injections, and
 - (b) for every cocone of injections, the factorizing morphism is an injection.

If I is a generator, then surjections are epimorphisms and (1) follows from the fact that every locally presentable category is co-wellpowered (see Adámek and Rosický (1994, Theorem 1.58)).

Examples 3.7. (1) Let **Pos** be the category of posets and monotone mappings. **Pos** is cartesian closed and I is the one-element poset 1. Surjections are surjective monotone mappings, i.e., epimorphisms. Injections are embeddings and $(Surj, Inj)$ is an ω -convenient factorization system.

(2) Let **Met** be the category of (generalized) metric spaces (i.e., with distances ∞ allowed) and nonexpanding maps. **Met** is a symmetric monoidal closed category where I is the one-element metric space 1 and $A \otimes B$ has the underlying set $A \times B$ and the metric

$$d((a, b), (a', b')) = d(a, a') + d(b, b').$$

Surjections are surjective nonexpanding mappings, injections are isometries and

$$(Surj, Inj)$$

is a factorization system (cf. Adámek and Rosický (2022, Example 3.16(1))). This factorization system is ω -convenient (see Adámek and Rosický (2022, Remark 2.5(2))).

(3) Let **Gra** be the cartesian closed category of graphs (i.e., sets with a symmetric binary relation) and graph homomorphisms. Then I is the one-element graph with a loop. Let V be a point (the one-element graph without a loop) and E an edge (the two-element graph without loops and with one edge). Then I is the coequalizer of two morphisms $V \rightarrow E$. Since $f : E \rightarrow I$ is not surjective, 3.5(1) implies that $(Surj, Inj)$ is not a factorization system on **V**.

4. Discrete Theories

Assumption 4.1. *Throughout the rest of the paper, we assume that \mathcal{V} is locally λ -presentable as a symmetric monoidal closed category and that the functor $\mathcal{V}_0(I, -) : \mathcal{V}_0 \rightarrow \mathbf{Set}$ preserves λ -presentable objects.*

Then the object X_0 of \mathcal{V} is λ -presentable whenever X is λ -presentable.

Definition 4.2. We say that a λ -ary \mathcal{V} -theory \mathcal{T} is *discrete* if every morphism $J1 \rightarrow JX$ of \mathcal{T} is a composition $J(\delta_X)f$ for some morphism $f : J1 \rightarrow JX_0$.

This means that operations of arity X are induced by those of arity X_0 .

Examples 4.3. (1) **Pos** is locally finitely presentable as a cartesian closed category. Finitely presentable posets are finite posets; hence, $\mathbf{Pos}(I, -) : \mathbf{Pos} \rightarrow \mathbf{Set}$ preserves finitely presentable objects. Let \mathcal{T} be a discrete finitary theory. Let Σ be the finitary signature whose n -ary operation symbols are \mathcal{T} -morphisms $J1 \rightarrow Jn$ where n denotes the discrete poset with n elements. Morphisms $J1 \rightarrow JX$ are X -ary terms and, since \mathcal{T} is discrete, they are the restrictions of usual X_0 -ary operation symbols on X (in the sense of 4.2). Pairs $f \leq g$, where $f, g : J1 \rightarrow JX$ yield inequations in context X in the sense of Adámek et al. (2021). Hence, \mathcal{T} yields a set of inequations in

context. The meaning of $f \leq g$ on a \mathcal{T} -algebra A is that $f_A(a_1, \dots, a_n) \leq g_A(a_1, \dots, a_n)$ provided that $a : n \rightarrow A$ factorizes through $X \rightarrow A$; here, $n = X_0$. Hence, \mathcal{T} -algebras coincide with algebras in the sense of Adámek et al. (2021). An example of a discrete finitary theory is the theory of ordered monoids which are commutative for comparable elements, i.e., $x \cdot y = y \cdot x$ for $x \leq y$.

Conversely, given a finitary signature Σ (with discrete arities), the free Σ -algebra FX on a poset X consists of X_0 -ary terms where $t \leq t'$ if $t(x_1, \dots, x_n)$ and $t'(x'_1, \dots, x'_n)$ have the same shape (i.e., we can get t' from t by changing variables x_1, \dots, x_n to x'_1, \dots, x'_n) and $x_i \leq x'_i$ in X for $i = 1, \dots, n$. Let E be a set of inequations in context. Following Adámek et al. (2021, 3.22), free E -algebra JX on a poset X is a quotient $q_X : FX \rightarrow JX$ of FX . If X is finite, every morphism $t : J1 \rightarrow JX$ factorizes through q_X as $t = q_X \bar{t}$. Hence, \bar{t} is an X_0 -ary term and $q_{X_0}(\bar{t})$ yields a factorization of t through $J(\delta_X)$. Thus, J is a discrete \mathcal{V} -theory.

(2) **Met** is locally \aleph_1 -presentable as a symmetric monoidal closed category. Since \aleph_1 -presentable objects are the metric spaces having cardinality $< \aleph_1$, **Met** $(I, -)$ preserves \aleph_1 -presentable objects. Let \mathcal{T} be a discrete \aleph_1 -ary theory. Like in (1), let Σ be the signature whose n -ary operation symbols are \mathcal{T} -morphisms $J1 \rightarrow Jn$ where n denotes the discrete metric space with n elements where $n \leq \omega$. Morphisms $J1 \rightarrow JX$ are X -ary terms and, since \mathcal{T} is discrete, they are the restrictions of usual X_0 -ary operation symbols on X . Pairs $f, g : J1 \rightarrow JX$ where $f, g : J1 \rightarrow JX$ and $d(f, g) \leq \varepsilon$ yield ω_1 -basic quantitative equations (Mardare et al., 2017). Hence, \mathcal{T} yields an ω_1 -basic quantitative equational theory of Mardare et al. (2017) (= a set of ω_1 -basic quantitative equations). The meaning $d(f, g) \leq \varepsilon$ on a \mathcal{T} -algebra A is that $d(f_A(a_1, \dots, a_n), g_A(a_1, \dots, a_n)) \leq \varepsilon$ provided that $a : n \rightarrow A$ factorizes through $X \rightarrow A$; here $n = X_0$. Hence, \mathcal{T} -algebras coincide with algebras in the sense of Mardare et al. (2017).

Conversely, every ω_1 -basic quantitative equational theory E (in a signature Σ with discrete arities) yields a discrete \aleph_1 -ary enriched theory \mathcal{T} . Analogously to (1), the free Σ -algebra on a countable metric space X consists X_0 -ary terms and $d(t, t') \leq \varepsilon$ if $t(x_1, \dots, x_n)$ and $t'(x'_1, \dots, x'_n)$ have the same shape and $d(x_i, x'_i) \leq \varepsilon$ in X for $i = 1, \dots, n$ (see Mardare et al. (2016)). This means that $d(t, t') = \infty$ if t and t' do not have the same shape. Following Mardare et al. (2016, 6.1), free E -algebra JX on a poset X is a quotient $q_X : FX \rightarrow JX$ of FX . The rest is the same as in (1).

Lemma 4.4. *Let \mathcal{T} be a λ -ary \mathcal{V} -theory. The following conditions are equivalent:*

- (1) \mathcal{T} is discrete,
- (2) every morphism $JY \rightarrow JX$ of \mathcal{T} with Y discrete is a composition $J(\delta_X)f$ with $f : JY \rightarrow JX_0$ and
- (3) for every morphism $g : JY \rightarrow JX$ of \mathcal{T} , there is a morphism $f : JY_0 \rightarrow JX_0$ such that $gJ(\delta_Y) = J(\delta_X)f$.

Proof. Clearly (3) \rightarrow (2) \rightarrow (1).

(1) \rightarrow (2). Every discrete object is a coproduct of copies of I . Since \mathcal{T} is an \mathcal{V}_λ -theory, J is given by free \mathcal{T} -algebras and thus it preserves coproducts. Hence, (1) implies (2).

(2) \rightarrow (3). It suffices to apply (2) on the composition $gJ(\delta_Y)$. □

Theorem 4.5. *A λ -ary \mathcal{V} -theory \mathcal{T} is discrete if and only if its induced monad preserves surjections.*

Proof. I. Let \mathcal{T} be discrete, $U : \mathbf{Alg}(\mathcal{T}) \rightarrow \mathcal{V}$ be the forgetful \mathcal{V} -functor and F its left \mathcal{V} -adjoint. Then the induced \mathcal{V} -monad is $T = UF$. Consider a λ -presentable object X , $\delta_X : X_0 \rightarrow X$ and $a : I \rightarrow UFX$. Since \mathcal{T} is discrete, the adjoint transpose $\tilde{a} : FI \rightarrow FX$ factorizes through $F\delta_X$, i.e., $\tilde{a} = F(\delta_X)b$ where $b : FI \rightarrow FX_0$. Hence, $a = UF(\delta_X)b$. We have proved that $UF(\delta_X)$ is surjective.

Consider an arbitrary X in \mathcal{V} and express it as a λ -directed colimit of λ -presentable objects X_m ($m \in M$). Since I is λ -presentable, $X_0 = \text{colim } X_{m0}$ and $\delta_X = \text{colim } \delta_{X_m}$. Hence, $UF(\delta_X) = \text{colim } UF(\delta_{X_m})$ and, since surjections are closed under λ -directed colimits (see 3.1), $UF(\delta_X)$ is a surjection.

Consider a surjective morphism $f : Y \rightarrow X$ in \mathcal{V} and $f_0 : Y_0 \rightarrow X_0$ be its underlying morphism. Then $f\delta_Y = \delta_X f_0$ and thus $UF(f)UF(\delta_Y) = UF(\delta_X)UF(f_0)$. Since epimorphisms in **Set** split, f_0 is a split epimorphism and $UF(f_0)$ is surjective. Hence, $UF(f)$ is surjective.

II. Let T be a λ -ary \mathcal{V} -monad on \mathcal{V} preserving surjections. Let \mathcal{T} be the corresponding λ -ary \mathcal{V} -theory. We will show that \mathcal{T} is discrete. Consider $f : FI \rightarrow FX$ and $\tilde{f} : I \rightarrow UFX$ its adjoint transpose. Since δ_X is surjective and T preserves surjections, $UF(\delta_X)$ is surjective. Thus, there is $g : I \rightarrow UFX_0$ such that $UF(\delta_X)g = \tilde{f}$. Let $\tilde{g} : FI \rightarrow FX_0$ be the adjoint transpose of g . Then $f = F(\delta_X)\tilde{g}$ and thus \mathcal{T} is discrete. □

Examples 4.6. (1) 4.5 (together with 4.3) gives the result of J. Adámek that sets of inequations in context in finitary signatures correspond to enriched finitary monads preserving surjections (see his talk “Finitary monads on **Pos**” at the conference Category theory CT20-21).

(2) Similarly, 4.5 gives the result of Adámek (2022) that ω_1 -basic quantitative equational theories (in signatures with discrete arities) correspond to enriched ω_1 -ary monads preserving surjections.

Remark 4.7. Let \mathcal{D}_λ consist of discrete λ -presentable objects. Then \mathcal{D}_λ -theories correspond to the λ -ary discrete Lawvere theories of Power (2005). Over **Pos** these theories are discrete in the sense of 4.2 and inequations of terms are in discrete contexts. Following Adámek et al. (2022) \mathcal{D}_ω -theories correspond to finitary enriched monads preserving reflexive comultiplicators. Over **Met**, \mathcal{D}_{\aleph_1} -theories correspond to unconditional quantitative theories of Mardare et al. (2017) and Mardare et al. (2016) and are discrete in the sense of 4.2. In general, we do not know whether discrete Lawvere theories are discrete in the sense of 4.2.

5. Birkhoff Subcategories

Assumption 5.1. Throughout this section, we assume, in addition to 4.1, that \mathcal{T} is a discrete λ -ary \mathcal{V} -theory and $(\text{Surj}, \text{Inj})$ is a proper μ -convenient factorization system in \mathcal{V}_0 where $\mu \leq \lambda$.

Observation 5.2. Recall that $(\text{Surj}, \text{Inj})$ is proper if surjections are epimorphisms and injections are monomorphisms. Let T be the \mathcal{V} -monad induced by \mathcal{T} . The underlying category $\mathbf{Alg}(T)_0$ is the category of algebras for the underlying monad T_0 . Let $U_0 : \mathbf{Alg}(T)_0 \rightarrow \mathcal{V}_0$ be the forgetful functor. We will say that a morphism $f : A \rightarrow B$ in $\mathbf{Alg}(T)_0$ is a surjection (injection) if $U(f)$ is a surjection (injection). Following Manes (1976, Chapter 3, Proposition 4.17), (surjections, injections) form a proper factorization system on $\mathbf{Alg}(T)_0$. Here we need that \mathcal{T} is discrete because then, following 4.5, T_0 preserves surjection. Given a morphism $f : A \rightarrow B$ of \mathcal{T} -algebras and

$$U_0 A \xrightarrow{e} C \xrightarrow{m} U_0 B$$

$(\text{Surj}, \text{Inj})$ -factorization of $U_0(f)$ then f factorizes as

$$A \xrightarrow{\bar{e}} \bar{C} \xrightarrow{\bar{m}} B$$

where $U_0 \bar{C} = C$, $U_0(\bar{e}) = e$ and $U_0(\bar{m}) = m$.

Since $\mathbf{Alg}(T)_0$ is locally presentable (strong epimorphisms, monomorphisms) is a factorization system on $\mathbf{Alg}(T)_0$ (see Adámek and Rosický (1994, 1.61)). Since injections are monomorphisms, strong epimorphisms in $\mathbf{Alg}(T)_0$ are surjections.

Following Manes (1976, Chapter 3, 3.1), a Birkhoff subcategory of $\mathbf{Alg}(T)_0$ is a full replete subcategory of $\mathbf{Alg}(T)_0$ closed under products, subalgebras and U_0 -split quotients. Here, a morphism $g : K \rightarrow L$ in $\mathbf{Alg}(T)_0$ is U_0 -split if $U_0(g)$ is a split epimorphism. Hence, for a Birkhoff subcategory \mathcal{L} of $\mathbf{Alg}(T)_0$, the reflections $\rho_K : K \rightarrow K^*$ are strong epimorphisms.

Lemma 5.3. *A Birkhoff subcategory of $\mathbf{Alg}(\mathcal{T})_0$ is \mathcal{V} -reflective in $\mathbf{Alg}(\mathcal{T})$.*

Proof. Let \mathcal{L} be a Birkhoff subcategory of $\mathbf{Alg}(\mathcal{T})$. Since \mathcal{L}_0 is reflective in $\mathbf{Alg}(\mathcal{T})_0$, it suffices to show that \mathcal{L} is closed in $\mathbf{Alg}(\mathcal{T})$ under cotensors (see Borceux (1994, 6.7.6). Consider V in \mathcal{V} and A in \mathcal{L} . Since the cotensor functor $[V, -] : \mathbf{Alg}(\mathcal{T}) \rightarrow \mathbf{Alg}(\mathcal{T})$ has a left \mathcal{V} -adjoint $V \otimes -$ (see Borceux (1994, 6.5.6)), and $\delta_V : V_0 \rightarrow V$ is an epimorphism, $[\delta_V, A] : [V, A] \rightarrow [V_0, A]$ is a monomorphism. Since $[V_0, A]$ is a power of A , $[V, A]$ is in \mathcal{L} . □

Hence, Birkhoff subcategories of $\mathbf{Alg}(\mathcal{T})_0$ are \mathcal{V} -subcategories and, in what follows, we will take them as subcategories of $\mathbf{Alg}(\mathcal{T})$. The following definition goes back to Hatcher (1970), Herrlich and Ringel (1972).

Definition 5.4. An \mathcal{T} -equation $p = q$ is a pair of morphisms $p, q : FY \rightarrow FX$ in $\mathbf{Alg}(\mathcal{T})$. A \mathcal{T} -equational theory E is a set of \mathcal{T} -equations.

A \mathcal{T} -algebra A satisfies a \mathcal{T} -equation $p = q$ if $hp = hq$ for every morphism $h : FX \rightarrow A$. It satisfies a \mathcal{T} -equational theory E if it satisfies all equations of E .

$\mathbf{Alg}(E)$ will be the full subcategory of $\mathbf{Alg}(\mathcal{T})_0$ consisting of \mathcal{T} -algebras satisfying all equations from E .

Proposition 5.5. *$\mathbf{Alg}(E)$ is a Birkhoff subcategory of $\mathbf{Alg}(\mathcal{T})$ for every \mathcal{T} -equational theory E .*

Proof. $\mathbf{Alg}(E)$ is clearly closed under products and subalgebras. Let $f : A \rightarrow B$ be a U_0 -split quotient of an algebra A satisfying E , i.e., there is $s : UB \rightarrow UA$ such that $U(f)s = \text{id}_{UB}$. Let $p, q : FY \rightarrow FX$ give an equation $p = q$ from E . Consider $h : FX \rightarrow B$ and $\tilde{h} : X \rightarrow UB$ be the adjoint transpose of h . Let $g = \tilde{s}\tilde{h}$ and $\tilde{g} : FX \rightarrow A$ be the adjoint transpose of g . Since $U(f)g = U(f)\tilde{s}\tilde{h} = \tilde{h}$, we have $f\tilde{g} = h$. Since $\tilde{g}p = \tilde{g}q$, we get

$$hp = f\tilde{g}p = f\tilde{g}q = hq.$$

Thus B satisfies E . □

Recall that an object V in \mathcal{V} is μ -generated with respect to Inj if $\mathcal{V}_0(V, -) : \mathcal{V}_0 \rightarrow \mathbf{Set}$ preserves μ -directed colimits of injections.

Definition 5.6. We say that a morphism $f : A \rightarrow B$ in \mathcal{V}_0 is a μ -pure epimorphism if it is projective with respect to μ -generated objects. Explicitly, for every μ -generated object X , all morphisms $X \rightarrow B$ factor through f .

Remark 5.7. (1) For $\mu = \lambda$, this concept was introduced in Adámek and Rosický (2004).

(2) Every μ -pure epimorphism $f : A \rightarrow B$ is an epimorphism. Indeed, assume that $uf = vf$ for $u, v : B \rightarrow C$. Consider a μ -generated object X and $g : X \rightarrow B$. Since g factors through f , we have $ug = vg$. Thus, $u = v$.

(3) Every split epimorphism is μ -pure.

(4) A morphism $f : A \rightarrow B$ ω -pure in \mathbf{Met} iff it is ω -reflexive in the sense of Mardare et al. (2017).

Definition 5.8. We say that \mathcal{L} is a μ -Birkhoff subcategory of $\mathbf{Alg}(\mathcal{T})$ if it is a Birkhoff subcategory closed under quotients $f : A \rightarrow B$ such that Uf is μ -pure.

Thus \mathcal{L} is a μ -Birkhoff subcategory of $\mathbf{Alg}(\mathcal{T})$ if and only if it is a full subcategory closed under products, subalgebras and μ -pure quotients.

Recall that \mathcal{V}_0 is *Inj*-locally μ -generated if it has a set \mathcal{X} of μ -generated objects with respect to *Inj* such that every object of \mathcal{V}_0 is a μ -directed colimit of objects from \mathcal{X} and injections (see Di Liberti and Rosický (2022)).

Remark 5.9. (1) Following Di Liberti and Rosický (2022, Remark 2.17), \mathcal{V}_0 is *Inj*-locally λ -generated.

(2) **Met** is *Inj*-locally ω -generated because finite metric spaces are ω -generated with respect to *Inj* (following Adámek and Rosický (2022, Remark 2.5(2))).

The following definition and theorem were motivated by Milius and Urbat (2019).

Definition 5.10. A \mathcal{T} -equation $p = q$, where $p, q : FY \rightarrow FX$, will be called μ -clustered if X is a coproduct of μ -generated objects.

Theorem 5.11. Assume 5.1 and, moreover, let \mathcal{V}_0 be *Inj*-locally μ -generated. Then \mathcal{L} is a μ -Birkhoff subcategory of $\mathbf{Alg}(\mathcal{T})$ if and only if $\mathcal{L} = \mathbf{Alg}(E)$ where all equations from E are μ -clustered.

Proof. I. Necessity: Let E consist of μ -clustered equations. We have to show that $\mathbf{Alg}(E)$ is closed in $\mathbf{Alg}(\mathcal{T})$ under quotients $f : A \rightarrow B$ such that Uf is μ -pure. So, let $f : A \rightarrow B$ be such a homomorphism. Consider an equation $p = q$ from E where $p, q : FY \rightarrow FX$ and $X = \coprod_i X_i$ with X_i μ -generated. Let $h : FX \rightarrow B$ and $\tilde{h} : X \rightarrow UB$ be its adjoint transpose. If $u_i : X_i \rightarrow X$ are coproduct injections, then there are $v_i : X_i \rightarrow UA$ such that $U(f)v_i = \tilde{h}u_i$ for every i . We get $v : X \rightarrow UA$ such that $U(f)v = \tilde{h}$. Therefore,

$$hp = f\tilde{v}p = f\tilde{v}q = hq.$$

Hence, B satisfies E .

II. Sufficiency: Conversely, let \mathcal{L} be a μ -Birkhoff subcategory of $\mathbf{Alg}(\mathcal{T})$. Following 5.3, \mathcal{L} is \mathcal{V} -reflective in $\mathbf{Alg}(\mathcal{T})$. Moreover, following 5.2, the reflections $\rho_K : K \rightarrow K'$ are strong epimorphisms.

Let $U' : \mathcal{L} \rightarrow \mathcal{V}_0$ be the restriction of U_0 on \mathcal{L} and F' its left adjoint. Let E be given by pairs (p, q) where $p, q : FY \rightarrow FX$ with $X = \coprod X_i$, where X_i are μ -generated with respect to *Inj* X_i , such that every $A \in \mathcal{L}$ satisfies the equation $p = q$. This is equivalent to $\rho_{FX}p = \rho_{FX}q$. Clearly, $\mathcal{L} \subseteq \mathbf{Alg}(E)$. Following I, $\mathbf{Alg}(E)$ is a μ -Birkhoff subcategory of $\mathbf{Alg}(\mathcal{T})$. Hence, $\mathbf{Alg}(E)$ is \mathcal{V} -reflective in $\mathbf{Alg}(\mathcal{T})$ and the reflections $\rho'_K : K \rightarrow K''$ are strong epimorphisms.

Let $U'' : \mathbf{Alg}(E) \rightarrow \mathcal{V}_0$ be the restriction of U on $\mathbf{Alg}(E)$ and F'' its left adjoint. We get $\rho' : F \rightarrow F''$ and $\tau : F'' \rightarrow F'$ such that $\tau\rho' = \rho_F$. Since both \mathcal{L} and $\mathbf{Alg}(E)$ are monadic, it suffices to prove that τ is a natural isomorphism.

Consider an arbitrary X in \mathcal{V} and express UFX as a μ -directed colimit

$$(z_m : Z_m \rightarrow UFX)_{m \in M}$$

of injections between objects Z_m which are μ -generated with respect to *Inj*. Let $\eta_Z : Z \rightarrow UFZ, Z$ in \mathcal{V} , be the adjunction units and $\tilde{z}_m : FZ_m \rightarrow FX$ be given as $U(\tilde{z}_m)\eta_{Z_m} = z_m$. Let

$$t_X : F(\coprod Z_m) \rightarrow FX$$

be the induced morphism, i.e., $t_X F(u_m) = \tilde{z}_m$ where $(u_m : Z_m \rightarrow \coprod Z_m)_{m \in M}$ is the coproduct. Then, $U(t_X)$ is a μ -pure epimorphism. Indeed, every morphism $f : Z \rightarrow UFX$ with Z μ -generated with respect to *Inj* factors through z_m for some $m \in M$ as $f = z_m g$ where $g : Z \rightarrow Z_m$. Hence,

$$\tilde{f} = \tilde{z}_m F(g) = t_X F(u_m) F(g)$$

and thus

$$f = U(\tilde{f})\eta_Z = U(t_X)UF(u_m g)\eta_Z.$$

We have to prove that τ_X is an isomorphism for every X . Consider the commutative diagram

$$\begin{CD} FU F'' X @>\varepsilon_{F'' X}>> F'' X \\ @V t_{UF'' X} VV @VV s V \\ F \coprod Z_m @>\rho'_{\coprod Z_m}>> F'' \coprod Z_m \end{CD}$$

where $\varepsilon : FU \rightarrow \text{Id}$ is the counit of the adjunction $F \dashv U$. The arrow s is given by $\rho'_{\coprod Z_m}$ being the reflection of $F \coprod Z_m$ to $\mathbf{Alg}(E)$. Since $U(\varepsilon_{F'' X})$ is a split epimorphism and $U(t_{UF'' X})$ is a μ -pure epimorphism, the composition $U(\varepsilon_{F'' X} t_{UF'' X})$ is a μ -pure epimorphism. Thus, $U(s)$ is a μ -pure epimorphism. Thus, it suffices to show that $F'' \coprod Z_m$ belongs to \mathcal{L} , i.e., that τ_Z is an isomorphism where $Z = \coprod Z_m$.

Since ρ_Z is a strong epimorphism, τ_Z is a strong epimorphism. Thus, it suffices to show that τ_Z is a monomorphism. Consider $p, q : FY \rightarrow F''Z$ such that $\tau_Z p = \tau_Z q$. We have to show that $p = q$. Consider the pullback

$$\begin{CD} FZ \times FZ @>\rho'_Z \times \rho'_Z>> F''Z \times F''Z \\ @V r VV @VV (p,q) V \\ P @>e>> FY \end{CD}$$

and take the compositions $p', q' : P \rightarrow FZ$ of the product projections with r . Let $p'', q'' : FUP \rightarrow FZ$ be the compositions of p' and q' with $\varepsilon_P : FUP \rightarrow P$. Since

$$\rho_Z p'' = \tau_Z \rho'_Z p'' = \tau_Z p e \varepsilon_P = \tau_Z q e \varepsilon_P = \rho_Z q'',$$

$p'' = q''$ belongs to E . Thus, we have $\rho'_Z p'' = \rho'_Z q''$, i.e., $p e = q e$. Since ρ'_Z is a strong epimorphism, it is a surjection (see 5.2). Following 3.2, e is a surjection, hence an epimorphism. Thus, $p = q$ and we have proved that τ_Z is a monomorphism. □

Definition 5.12. We will call a category \mathcal{K} *strongly connected* if for every pair of objects K and K' of \mathcal{K} , where K' is not initial, there is a morphism $K \rightarrow K'$.

Lemma 5.13. *The following conditions are equivalent:*

- (1) \mathcal{K} is strongly connected, and
- (2) for a coproduct $\coprod_{m \in M} K_m$ of non-initial objects K_m , coproduct components $u_m : K_m \rightarrow \coprod_{m \in M} K_m$ are split monomorphisms.

Proof. If \mathcal{K} is strongly connected and $u_m : K_m \rightarrow \coprod_{m \in M} K_m$ a coproduct of non-initial objects, then there is a cocone $v_n : K_m \rightarrow K_n$ for every n with $v_n = \text{id}_{K_n}$. Hence u_n is a split monomorphism.

Conversely, (2) applied to $K \coprod K'$ yields $K \rightarrow K'$. □

Remark 5.14. In 5.13, all subcoproduct morphisms $\coprod_{m \in N} X_m \rightarrow \coprod_{m \in M} X_m$, $N \subseteq M$, are split monomorphisms. In fact, $\coprod_{m \in M} X_m = (\coprod_{m \in N} X_m) \coprod (\coprod_{m \notin N} X_m)$.

Theorem 5.15. Assume 5.1 and, moreover, let \mathcal{V}_0 be *Inj*-locally μ -generated, strongly connected and let U preserve μ -directed colimits of injections. Then \mathcal{L} is a μ -Birkhoff subcategory of $\mathbf{Alg}(\mathcal{T})$ if and only if $\mathcal{L} = \mathbf{Alg}(E)$ where all equations $p = q$ from E have $p, q : FY \rightarrow FX$ with X and Y being μ -generated with respect to *Inj*.

Proof. Sufficiency is I. of the proof of 5.11. Conversely, let \mathcal{L} be a μ -Birkhoff subcategory of $\mathbf{Alg}(\mathcal{T})$. Following 5.11, $\mathcal{L} = \mathbf{Alg}(E)$ where all equations $p = q$ from E have $p, q : FY \rightarrow FX$ with X being a coproduct of objects μ -generated with respect to *Inj*. Express Y as a μ -directed colimit $(y_m : Y_m \rightarrow Y)_{m \in M}$ of objects Y_m μ -generated with respect to *Inj*. Then $F(y_m) : FY_m \rightarrow FY$ is a μ -directed colimit and $\rho_{FX}p = \rho_{FX}q$ if and only if $\rho_{FX}pF(y_m) = \rho_{FX}qF(y_m)$ for every $m \in M$. Thus we have reduced our equations to $p = q$ where $p, q : FY \rightarrow FX$ with X being a coproduct of objects μ -generated with respect to *Inj* and Y being μ -generated with respect to *Inj*.

Express $X = \coprod_{m \in M} X_m$ as a μ -directed colimit of subcoproducts $X_N = \coprod_{m \in N} X_m$ where $|N| < \mu$ and X_m are non-initial. Following 5.14, X is a μ -directed colimit $x_N : X_N \rightarrow X$ of split monomorphisms $x_{NN'} : X_N \rightarrow X_{N'}$ where $N \subseteq N'$. Hence, UFx is a μ -directed colimit $UF(x_N) : UFX_N \rightarrow UFX$ of split monomorphisms $UF(x_{NN'}) : UFX_N \rightarrow UFX_{N'}$. Following 3.5, $UF(x_{NN'})$ are injections. Let $\tilde{p}, \tilde{q} : Y \rightarrow UFX$ be the adjoint transposes of $p, q : FY \rightarrow FX$. Since Y is μ -generated with respect to *Inj*, there is $N \subseteq M, |N| < \mu$ and $p', q' : Y \rightarrow UFX_N$ such that $\tilde{p} = x_N p'$ and $\tilde{q} = x_N q'$.

Now, $p = q$ is an equation from E , iff $U(\rho_{FX})\tilde{p} = U(\rho_{FX})\tilde{q}$. Since x_N is a split monomorphism, this is equivalent to $U(\rho_{FX_N})p' = U(\rho_{FX_N})q'$, hence to $\rho_{X_N}p^* = \rho_{X_N}q^*$ where $p^*, q^* : FY \rightarrow FX_N$ are adjoint transposes of p', q' . Since X_N is μ -generated with respect to *Inj* (see Di Liberti and Rosický (2022, Lemma 2.13)), the proof is finished. \square

Remark 5.16. Moreover, $\mathbf{Alg}(E)$ is closed in $\mathbf{Alg}(\mathcal{T})$ under μ -directed colimits of injections.

Indeed, let $k_m : K_m \rightarrow K$ be a μ -directed colimit of injections where K_m satisfy E . Consider an equation $p = q$ from E where $p, q : FY \rightarrow FX$. Since U preserves μ -directed colimits of injections, F preserves μ -generated objects (see Di Liberti and Rosický (2022, Lemma 3.11)). Hence, a morphism $h : FX \rightarrow K$ factors through some $k_m : K_m \rightarrow K, h = k_m h'$. Since K_m satisfies $E, h'p = h'q$. Thus, $hp = hq$ and K satisfies E .

Remark 5.17. The assumption that \mathcal{K} is strongly connected is needed in 5.15 because 5.16 is not valid for the category $\mathbf{Set}^{\mathbb{N}}$ of \mathbb{N} -sorted sets. Indeed, let T be a finitary monad in $\mathbf{Set}^{\mathbb{N}}$. Since epimorphisms in $\mathbf{Set}^{\mathbb{N}}$ split, every Birkhoff subcategory of $(\mathbf{Set}^{\mathbb{N}})^T$ is ω -Birkhoff. Following Adámek et al. (2012), it does not need to be closed under directed colimits. (Take the full subcategory of $\mathbf{Set}^{\mathbb{N}}$ consisting of all $(X_n)_{n \in \mathbb{N}}$ such that either $X_n = \emptyset$ for some n or $X_n = 1$ for every n .)

As a consequence of 5.15, we get the Birkhoff theorems over **Met** from Mardare et al. (2017), see Adámek (2022, Theorem 2.14). Here, under an ω -ary theory over **Met** we mean an \mathcal{F} -theory where \mathcal{F} consists of finite metric spaces, i.e., of finitely generated metric spaces with respect to *Inj*.

Manes (1976) defined a *Surj-Birkhoff subcategory* \mathcal{L} of $\mathbf{Alg}(\mathcal{T})_0$ as a full replete subcategory closed under products, *Inj*-subalgebras and U_0 -split quotients. Here, a morphism $g : K \rightarrow L$ in $\mathbf{Alg}(\mathcal{T})_0$ is an *Inj*-subalgebra of L if g is an injection. Since the monad T preserves surjections, Manes (1976, Chapter 3, Theorem 4.23) shows that a full subcategory \mathcal{L} of $\mathbf{Alg}(\mathcal{T})_0$ is a *Surj-Birkhoff subcategory* if and only if it is a reflective subcategory such that $U_0(\rho_K)$ is a surjection for every reflection $\rho_K : K \rightarrow K^*$. Following Manes (1976, Chapter 3, Theorem 3.3), \mathcal{L} is monadic over \mathcal{V}_0 . Hence, every Birkhoff subcategory of $\mathbf{Alg}(\mathcal{T})_0$ is a *Surj-Birkhoff subcategory*.

Corollary 5.18. Let Σ be a μ -ary (discrete) signature where $\mu = \omega$ or $\mu = \omega_1$. Then a class of quantitative Σ -algebras is a μ -*Surj-Birkhoff subcategory* of $\mathbf{Alg}(\Sigma)$ if and only if it is given by an μ -basic quantitative equational theory.

Proof. At first, we observe that, for a discrete μ -ary **Met**-theory \mathcal{T} , the functor $U: \mathbf{Alg}(\mathcal{T}) \rightarrow \mathbf{Met}$ preserves μ -directed colimits of isometries. For $\mu = \omega_1$ it is evident and, for $\mu = \omega$, it follows from Rosický (2021, Theorem 3.19).

Following 4.3(2), $\mathbf{Alg}(\Sigma) \cong \mathbf{Alg}(\mathcal{T})$ for a discrete μ -ary **Met**-theory \mathcal{T} . Given a μ -basic quantitative equational theory E then, following 4.3(2) and I. of the proof of 5.11, $\mathbf{Alg}(E)$ is a μ -Surj-Birkhoff subcategory of $\mathbf{Alg}(\Sigma)$.

Conversely, let \mathcal{L} be a μ -Surj-Birkhoff subcategory of $\mathbf{Alg}(\Sigma)$. We extend Σ to a new signature Σ^* by adding, for every $\varepsilon > 0$, binary operations f_ε and g_ε satisfying the quantitative equations

- (1) $f_\varepsilon(x, y) =_\varepsilon g_\varepsilon(x, y)$ and
- (2) $x =_\varepsilon y \vdash f_\varepsilon(x, y) = x, \quad x =_\varepsilon y \vdash g_\varepsilon(x, y) = y.$

Then subalgebras of Σ^* -algebras satisfying (1) and (2) are *Inj*-subalgebras. Indeed, let B be a subalgebra of a Σ^* -algebra A satisfying (1) and (2) and consider $a, b \in B$ such that $d_B(a, b) > d_A(a, b)$. Following (2) for $\varepsilon = d_A(a, b)$, we have

$$(f_\varepsilon)_B(a, b) = (f_\varepsilon)_A(a, b) = a, \quad (g_\varepsilon)_B(a, b) = (g_\varepsilon)_A(a, b) = b,$$

which contradicts (1).

Let \mathcal{L}^* consist of all Σ^* -algebras satisfying (1) and (2) whose Σ -reduct is in \mathcal{L} . Then \mathcal{L}^* is a μ -Birkhoff subcategory of $\mathbf{Alg}(\Sigma^*)$. Let \mathcal{T}^* be the discrete μ -ary **Met**-theory given by Σ^* , (1) and (2). Following 5.15, $\mathcal{L}^* = \mathbf{Alg}(E^*)$ where equations $p = q$ from E^* have $p, q: F^*Y \rightarrow F^*X$ with X and Y being μ -generated with respect to *Inj*. Let E' consist of those equations from E^* which do not contain the added operations f_ε and g_ε . Clearly, $\mathbf{Alg}(E') \subseteq \mathcal{L}$. Consider $A \in \mathcal{L}$ and interpret f_ε and g_ε on A as follows: $(f_\varepsilon)_A(x, y) = x$ and $(g_\varepsilon)_A(x, y) = y$ if $d(x, y) = \varepsilon$ and $(f_\varepsilon)_A = (g_\varepsilon)_A$ otherwise. Since we get an E^* -algebra, $\mathcal{L} \subseteq \mathbf{Alg}(E')$. \square

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Appendix A. Appendix

Rosický (2021, Proposition 5.1) claims that finite products commute with reflexive coequalizers in **Met**. Equivalently, that the functor

$$X \times - : \mathbf{Met} \rightarrow \mathbf{Met}$$

preserves reflexive coequalizers. Since this functor preserves coproducts, the preservation of reflexive coequalizers is equivalent to the preservation of all colimits, i.e., following the Special Adjoint Functor Theorem, to the cartesian closedness of **Met**. Thus Rosický (2021, Proposition 5.1) is not true and I am grateful to Jason Parker for pointing this out.

In fact, Rosický (2021, Proposition 5.1) proves that the functor

$$X \times - : \mathbf{Dist} \rightarrow \mathbf{Dist}$$

preserves reflexive coequalizers where **Dist** is the category of distance spaces and nonexpanding mappings. Recall that a *distance space* is equipped with a metric $d : X \rightarrow [0, \infty]$ satisfying $d(x, y) = d(y, x)$ and $d(x, x) = 0$. The category **Dist** is cartesian closed (Adámek and Reiterman, 1990).

The category **PMet** of pseudometric spaces (i.e., with the triangle inequality added) is a reflective subcategory of **Dist**; the reflection of (X, d) is obtained by the *pseudometric modification* d^* of d :

$$d^*(x, z) = \inf\left\{\sum_{i=0}^{n-1} d(y_i, y_{i+1}) \mid n \geq 1, y_i \in X, y_0 = x, y_n = z\right\}.$$

Hence, **Met** is a reflective subcategory of **Dist**. We have just explained that the reflector $F : \mathbf{Dist} \rightarrow \mathbf{Met}$ cannot preserve finite products. We will show that it does not even preserve finite powers.

Example A.1. Consider the distance spaces $A_0 = \{x, y, z\}$ and $A_1 = \{u, v, w\}$ where $d(x, y) = d(v, w) = 2$, $d(y, z) = d(u, v) = 1$ and $d(x, z) = d(u, w) = \infty$. Let A be the coproduct $A_0 + A_1$. Then the metric modification FA only changes $d(x, z)$ and $d(u, w)$ to 3. Hence, the distance of $[x, u]$ and $[z, w]$ in $FA \times FA$ is 3.

In $A \times A$, we have $(d[x, u], [y, v]) = d([y, v], [z, w]) = 2$, hence $[d[x, u], [z, w]] = 4$ in $F(A \times A)$.

The distance space A_0 above is the quotient in **Dist** of the coproduct $A_{00} + A_{01} + A_{02}$ of metric spaces $A_{00} = \{x, y\}$, $A_{01} = \{y', z\}$ and $A_{02} = \{x', z'\}$ where $d(x, y) = 2$, $d(y', z) = 1$ and $d(x', z') = \infty$ modulo the equivalence relation $x \sim x'$, $y \sim y'$ and $z \sim z'$. Thus, FA_0 is the quotient of $A_{00} +$

$A_{01} + A_{02}$ modulo this equivalence relation in **Met**. Similarly, A_1 is the quotient in **Dist** of the coproduct $A_{10} + A_{11} + A_{12}$ of metric spaces $A_{10} = \{u, v\}$, $A_{11} = \{v', w\}$ and $A_{12} = \{u', w'\}$ modulo the equivalence relation $u \sim u'$, $v \sim v'$ and $w \sim w'$. Hence, FA_1 is the quotient of $A_{10} + A_{11} + A_{12}$ modulo this equivalence relation in **Met**. Consequently, FA is the quotient of the equivalence relation \sim on $B = A_{00} + A_{01} + A_{02} + A_{10} + A_{11} + A_{12}$ in **Met**. Similarly, $F(A \times A)$ is the quotient of $\sim \times \sim$ on $B \times B$.

Example A.1 thus shows that quotients of equivalence relations do not commute with finite powers in **Met**. Hence, reflexive coequalizers do not commute with finite powers in **Met**. Consequently, Rosický (2021, Corollary 5.2, Corollary 5.3, Lemma 5.4 and Example 5.6) are false.